Quadrangulations and 2-Colorations

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Abstract

Any metric quadrangulation (made by segments of straight line) of a point set in the plane determines a 2-coloration of the set, such that edges of the quadrangulation can only join points with different colors. In this work we focus in 2-colorations and study whether they admit a quadrangulation or not, and whether, given two quadrangulations of the same 2-coloration, it is possible to carry one into the other using some local operations, called *diagonal slides* and *diagonal rotation*. Although the answer is negative in general, we can show a very wide family of 2-colorations, called *onions 2-coloration*, that are quadrangulable and which graph of quadrangulations is always connected.

1 Introduction

Given a set S, either a polygon or a point set, a quadrangulation of S is a partition of the interior of S, if S is a polygon, or of the convex hull of S, if S is a point set, into quadrangles (quadrilaterals) obtained by inserting edges between pairs of points (diagonals between vertices of the polygon) such that the edges intersect each other only at their end points. Not all polygons or point sets admit quadrangulations, even when the quadrangles are not required to be convex. In the study of finite element methods and scattered data interpolation, it has recently been shown that quadrangulations of point sets may be more desirable objects than triangulations [2]. The quadrangulations of polygons have been also investigated in Computational Geometry, mostly in the context of guarding or illumination problems.

From now on we call a polygon or point set *quadrangulable* if it admits a quadrangulation without adding any additional point (Steiner point).

There are two different characterizations of quadrangulable point sets:

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- if and only if there exists a triangulation of the set such that its dual graph contains a perfect matching [5].

- if and only if it has an even number of points in its convex hull [1].

A quadrangulation is constructed (with the addition of a Steiner point to obtain an even number of points in the convex hull, if necessary) in $\Theta(n \log n)$.

For a more complete vision on quadrangulations we recommend Toussaint's survey [6].

Given a quadrangulation of a point set in the plane, it determines a 2-coloration of the set, such that edges of the quadrangulation can only join points with different colors. In this work we focus on 2-colorations and study whether they admits a quadrangulation or not (Section 2), and whether, given two quadrangulations of the same 2-coloration, it is possible to carry one into the other using some local operations (Section 3). Finally, in Section 4 we present a very wide family of 2-colorations, called *onions 2-coloration*, that are quadrangulable and which graph of quadrangulations is always connected.

2 2-colorations and quadrangulations

Suppose we have a 2-colorated point set S in the plane and we want to know if it is possible to construct a quadrangulation of its convex hull. A similar condition to the one given by [1] is, in this context, evident:

Lemma 1 *A necessary condition for a 2-coloration of a point set S to admit a quadrangulation (to be* quadrangulable*) is that*

1. *the number of points of the convex hull of S is even; and*

2. *consecutive points of the hull have different color.*

But even when the conditions of Lemma 1 are fulfilled, it is easy to find non-quadrangulable 2-colorations, as the one at the left of Figure 1. In order to construct a quadrangulation, point 1 cannot be joined with c because then b can be joined with no black point but 1. So we draw an edge from 1 to b. Now b cannot be joined either with 3 or with 4, because then c or 2, respectively, would be isolated, and cannot be part of any quadrangulation. But if

we join b with 2 the only way to complete a quadrilateral is to match 2 and d, that leaves 3 isolated. It is important to remark that we are talking about 2-colorations instead sets of points. Thus, while the 2-coloration at the left of Figure 1 is not quadrangulable, the underlying point set is, as we see in the right picture.
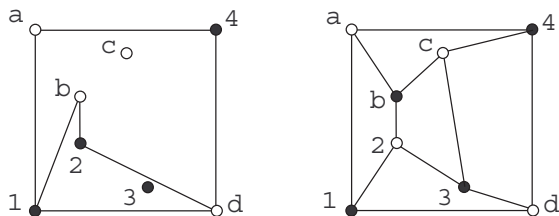

Figure 1: The set is either quadrangulable or not depending on the coloration.

Notice that in the right picture we have interchanged the colors of 2 and b, obtaining a set with two convex layers, both of them made by points with alternate colors. This is a interesting configuration since, as we will see in Section 4, it is always quadrangulable.
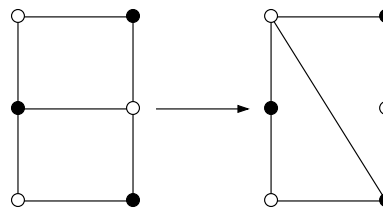
3 Diagonal transformation in quadrangulations

Nakamoto [4], working with topological quadrangulations on surfaces, defines two diagonal transformations; the diagonal slide and the diagonal rotation, that are shown in Figure 2. Note that, while the diagonal slide does not modify the coloration of the set, the diagonal rotation changes the color of the *center* of rotation (because, in other case, points with the same color are joined).

Since the same point set can have different colorations, it is not always possible to change any two quadrangulations one into each other using only diagonal slides. In Figure 3 two colorations of the same point set are shown; one with four black and four white points, and another with five white and three black points. Since diagonal slides preserve colorations, it is not possible to use them to transform one quadrangulation into the other. However, it is easy to see that this can be done using also diagonal rotations.

Nakamoto [4] proved that in any closed surface it is always possible to carry one topological quadrangulation of a set into any other if

1. both diagonal slides and rotations are allowed; or

2. both quadrangulations have the same number of points of each color by using only diagonal slides.

diagonal slide
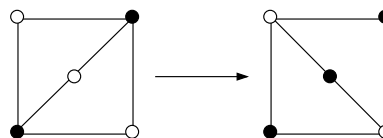
diagonal rotation

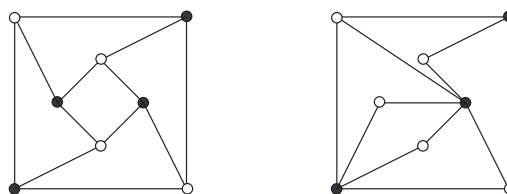Figure 2: Diagonal transformations on quadrangulations.

Figure 3: Quadrangulations of the same set with different colorations.

This can be seen in terms of the connectivity of the graph of quadrangulations. The *graph of quadrangulations of a point set* is the graph having all the quadrangulations of the set as nodes, and with adjacentcies corresponding to diagonal slides or diagonal rotations. Similarly, the *graph of quadrangulations of a 2-coloration* has as nodes the quadrangulations of a given 2-coloration. Since diagonal rotations change the 2-coloration, the adjacentcies are determined only by diagonal slides.

Both graph of quadrangulations are, in general, not connected. In Figure 4, it is shown a 2-coloration that admits only two quadrangulations, being not possible to perform any diagonal slide. If we also allow diagonal rotations it can be shown that it is not possible to transform one quadrangulation into the other. This gives rise to the following theorems:

Theorem 2 *There are 2-colorations with disconnected graph of quadrangulations.*

Theorem 3 *There are point sets with disconnected graph of quadrangulations.*
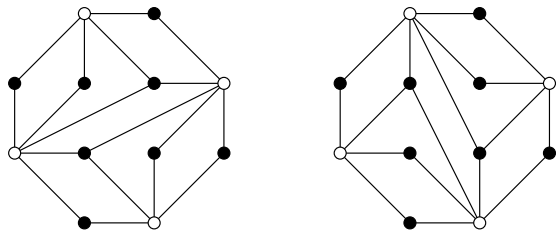
Figure 4: A set with disconnected graph of quadrangulations.

Figure 5: An onion 2-coloration and its onion layers.

Our example has disconnected graph of quadrangulations both as a 2-coloration and as a point set. An open problem is to determine if both things always come together, or if there exit point sets with connected graph of quadrangulations that admit 2-colorations which graph of quadrangulations is not.

In spite of the graph of quadrangulations of an arbitrary 2-coloration is, in general, not connected, in the next section we present a wide family of 2-colorations having this property.

4   Onion 2-colorations

If a set of sites have an even number of vertices in its convex hull, then it is quadrangulable, and vice-versa [1]. This is rewritten for 2-colorations in Lemma 1, but only as a necessary condition, since we find non-quadrangulable 2-colorations that fulfill it, as the one we saw in Figure 1. But, what about if we extend Lemma 1 to the interior of the set? If the convex hull of the 2-coloration fulfill the lemma, we remove it and examine the convex hull of the remaining points, and so on. A 2-coloration with this property is quadrangulable and its graph of quadrangulations is connected.

We call *onion 2-coloration* to a 2-coloration of a point set such that all its convex layers have an even number of points with alternate colors. An *onion layer* of an onion 2-coloration is the set of edges that are part of a convex layers of the set. We call O_i, with $i = 0, \ldots, l$, to the onion layers of the onion 2-coloration, such that O_i is inside the polygon defined by O_j if $i > j$ (Figure 5). Note that the polygon defined by O_l does not contain any point of the onion 2-coloration and that O_0, the convex hull, is always included in every quadrangulation of the onion 2-coloration. By definition, the points on every onion layer satisfy Lemma 1, that implies the following result.

Proposition 4 *Onion 2-colorations are quadrangulable.*

The main idea of the proof is to draw a triangulation joining points between two consecutive onion layers. By deleting the edges matching points with the
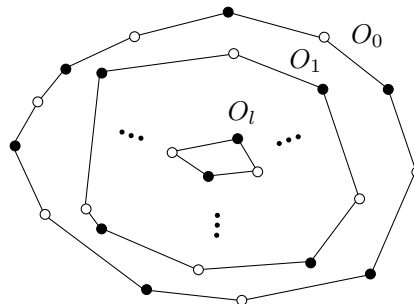
same color (Figure 6) we obtain a quadrangulation of the onion 2-coloration. It should be note that we are drawing quadrangulations of convex polygons with a convex hole, being the general case, decide whether a polygon with holes admits a quadrangulation, an NP-complete problem [3].
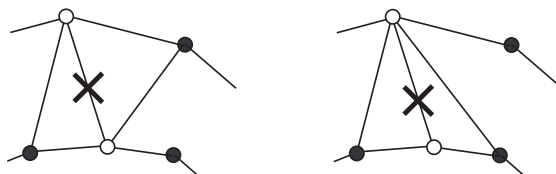

Figure 6: By deleting the diagonals between points with the same color we obtain quadrilaterals.

But onion 2-colorations are not the only quadrangulable 2-colorations, since there are quadrangulable 2-colorations with non alternate colors in some of its convex layer (Figure 7) or with a odd number of points on them.
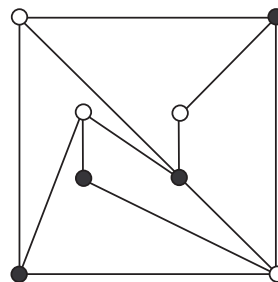

Figure 7: The colors of the inner convex layer are not alternate.

In addition to be quadrangulable, onion 2-colorations have connected graph of quadrangulations. The proof is based in the following lemmas:

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Lemma 5 *Given two quadrangulations of an onion 2-coloration containing all their onion layers, we can transform one into the other by only using diagonal slides.*

Lemma 6 *Any quadrangulation of an onion 2-coloration can be carried into another containing its onion layers using diagonal slides.*

From these lemmas it can be easily proved the connectivity of the graph of quadrangulations of any onion 2-coloration.

Theorem 7 *The graph of quadrangulations of an onion 2-coloration is non-empty and connected.*

In particular, if the onion 2-coloration have only one layer, we obtain the following result:

Corollary 8 *The graph of quadrangulations of any quadrangulable 2-coloration in convex position is connected.*

5 Conclusions and open problems

Two main ideas can be extracted from this work: to be quadrangulable depends on the 2-coloration of the set, and the graph of quadrangulations of both a 2-coloration and a point set is, in general, not connected. However, there exits a wide family of 2-colorations, the onion 2-colorations, that are quadrangulable and which graph of quadrangulations is connected.

There are several questions that appear all along the present work. One is to explore new conditions for a 2-coloration to be quadrangulable, searching for new families of quadrangulable 2-colorations. Related to the graph of quadrangulations, an interesting approach is to study the relationship, if it exits, between the connectivity of the graph and the coloration of the set. And, since the example presented (Figure 4) of a set with disconnected graph of quadrangulations have rows with until four collinear points, it would be convenient to construct a new example in general position. Probably this implies to work with sets with greater cardinal and complexity. Finally, another line for future works is, since they also admit 2-colorations, to extend this study from quadrangulations to 2n-lations of point sets.

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